12.

Abductive Reasoning: Explanation and Diagnosis

Outline

- Introduction to Abductive Reasoning
- Explanation & Diagnosis
- Computing Explanations
- Reading Material

• Deduction

 an analytic process based on the application of general rules to particular cases with the inference of a result

Induction

synthetic reasoning which infers the rule from the case and the result

• Abduction

 another form of synthetic inference but of the case from a rule and a result

- Probational adoption of a hypothesis as explanation for observed facts results according to known laws
- Weak form of inference because we cannot say that we believe in the truth of the explanation

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grass-is-wet ← rained-last-night
grass-is-wet ← sprinkler-was-on
shoes-are-wet ← grass-is-wet
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- Form of non-monotonic reasoning
- Multiple explanations might exist

Abductive Reasoning

Given KB, and an α that I do *not* believe,

what would be sufficient to make me believe that α was true?

- or what else would I have to believe for α to become an implicit belief?
- or what would explain α being true?

Deduction: given $(p \supset q)$, from p, deduce q

Abduction: given $(p \supset q)$, from q, abduce p

p is sufficient for q or one way for q to be true is for p to be true

Also induction: given $p(t_1), q(t_1), ..., p(t_n), q(t_n)$, induce $\forall x (p(x) \supset q(x))$

Can be used for causal reasoning: (*cause* \supset *effect*)

One simple version of diagnosis uses abductive reasoning

 KB has facts about symptoms and diseases including: (*Disease* ∧ *Hedges* ⊃ *Symptoms*)
 Goal: find disease(s) that best explain observed symptoms

Observe: we typically do not have knowledge of the form (Symptom $\land ... \supset$ Disease)

so reasoning is not deductive

Example:

(tennis-elbow \supset sore-elbow) (tennis-elbow \supset tennis-player) (arthritis \land untreated \supset sore-joints) (sore-joints \supset sore-elbow \land sore-hip) Explain: sore-elbow

. . .

Want: tennis-elbow, (arthritis ^ untreated),

Non-uniqueness: multiple equally good explanations

+ logical equivalences: (untreated ^ __arthritis)

Given KB, and β to be explained, we want an $\alpha\,$ such that

1. α is sufficient to account for β

 $\mathsf{KB} \cup \{\alpha\} \mid= \beta \quad \text{ or } \quad \mathsf{KB} \mid= (\alpha \supset \beta)$

2. α is not ruled out by KB

 $\mathsf{KB} \cup \{\alpha\} \text{ is consistent} \quad \text{or} \quad \mathsf{KB} \ | \neq \neg \alpha$

3. α is as simple as possible

parsimonious : as few *terms* as possible explanations should not unnecessarily strong or unnecessarily weak

4. α is in the appropriate vocabulary

atomic sentences of α should be drawn from **H**, possible hypotheses in terms of which explanations are to be phrased

e.g. diseases, original causes

Call such α an <u>explanation</u> of β wrt KB

otherwise $(p \land \neg p)$ would count as an explanation

- e.g. KB = { $(p \supset q), \neg r$ } and $\beta = q$ $\alpha = (p \land s \land \neg t)$ is too strong $\alpha = (p \lor r)$ is too weak
- e.g. sore-elbow explains sore-elbow trivial explanation
 - sore-joints explains sore-elbow may or may not be suitable

Example

- For the observation, shoes-are-wet,
 - the explanation grass-is-wet is not in the appropriate vocabulary
 - But it-rained-last-night, and sprinkler-was-on are in the appropriate vocabulary
 - {it-rained-last-night, sprinkler-was-on} is not a minimal explanation, but {it-rained-last-night} is a minimal explanation

- It is sufficient to consider
 - how to explain an atom or a literal
 - Conjunction of literals (or negation of a clause)

From criteria of previous slide, we can simplify explanations in the propositional case, as follows:

• To explain an arbitrary wff β , it is sufficient to choose a new letter p, add $(p \equiv \beta)$ to KB, and then explain p.

 $\mathsf{KB} \mid = (E \supset \beta) \quad \text{iff} \quad \mathsf{KB} \cup \{(p \equiv \beta)\} \mid = (E \supset p)$

• Any explanation will be (equivalent to) a conjunction of literals (that is, the negation of a clause)

Why? If α is a purported explanation, and DNF[α] = $(d_1 \lor d_2 \lor ... \lor d_n)$ then each d_i is also an explanation that is no less simple than α

A simplest explanation is then the negation of a clause with a *minimal* set of literals

So: to explain a literal ρ , it will be sufficient to find the minimal clauses *C* (in the desired vocabulary) such that

1. KB |=
$$(\neg C \supset \rho)$$
 or KB |= $(C \cup \{\rho\})$ sufficient

2. KB
$$\not\models C$$
 consistent

Prime implicates

A clause C is a prime implicate of a KB iff

- 1. KB |= C
- 2. For no $C^* \subset C$, KB |= C^*

Note: For any clause C, if KB |= C, then some subset of C is a prime implicate

Example: KB = { $(p \land q \land r \supset g), (\neg p \land q \supset g), (\neg q \land r \supset g)$ }

Prime implicates:



For explanations:

- want minimal C such that KB |= $(C \cup \{\rho\})$ and KB | $\neq C$
- so: find prime implicates *C* such that $\rho \in C$; then $\neg(C - \rho)$ must be an explanation for ρ

Example: explanations for g in example above

• 3 prime implicates contain g, so get 3 explanations: $(\neg p \land q)$, r, and g

Given KB, to compute explanations of literal ρ in vocabulary H:

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calculate the set \{\neg(C - \rho) \mid C \text{ is a prime implicate and } \rho \in C\}
prime implicates containing \rho
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But how to compute prime implicates?

Can prove: Resolution is complete for non-tautologous prime implicates

 $\mathsf{KB} \models C \text{ iff } \mathsf{KB} \rightarrow C \text{ completeness for [] is a special case!}$

So: assuming KB is in CNF, generate *all* resolvents in language **H**, and retain those containing ρ that are minimal

Could pre-compute all prime implicates, but there may be *exponentially* many, even for a Horn KB

explain: E_n

Circuit example

Components



Circuit behaviour

Truth tables for logical gates

and (0,0) = 0, and (0,1) = 0, ... or (0,0) = 0, or (0,1) = 1, ... xor (0,0) = 0, xor (0,1) = 1, ...

Normal behaviour

And gate(x) $\land \neg Ab(x) \supset out(x) = and(in1(x), in2(x))$ Orgate(x) $\land \neg Ab(x) \supset out(x) = or(in1(x), in2(x))$ Xorgate(x) $\land \neg Ab(x) \supset out(x) = xor(in1(x), in2(x))$

Abnormal behaviour: fault models

Examples

 $[Orgate(x) \lor Xorgate(x)] \land Ab(x) \supset out(x) = in2(x)$ (short circuit)

Other possibilities ...

- some abnormal behaviours may be inexplicable
- some may be compatible with normal behaviour on certain inputs

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Given KB as above + input settings
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e.g. KB \cup \{in1(f) = 1, in2(f) = 0, in3(f) = 1\}
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we want to explain observations at outputs

e.g. $(out1(f) = 1 \land out2(f) = 0)$

in the language of Ab

We want conjunction of Ab literals α such that KB \cup Settings $\cup \{\alpha\} \models Observations$

Compute by "propositionalizing":

For the above, x ranges over 5 components and u, v range over 0 and 1.

Easiest to do by preparing a table ranging over all $\rm Ab$ literals, and seeing which conjunctions entail the observations.



	Ab(b1)	Ab(b2)	Ab(a1)	Ab(a2)	Ab(01)	Entails observation?
1.	Y	Y	Y	Y	Y	N
2.	Y	Y	Y	Y	Ν	Ν
3.	Y	Y	Y	Ν	Y	Ν
4.	Y	Y	Y	Ν	N	Ν
5.	Y	Y	Ν	Y	Y	Y
6.	Ý	Ŷ	N	Ŷ	N	N
7.	Y	Y	Ν	Ν	Y	Y
8.	Y	Y	Ν	N	N	Y
9.	Y	Ν	Y	Y	Y	Ν
10.	Y	Ν	Y	Y	Ν	Ν
11.	Y	Ν	Y	Ν	Y	Ν
12.	Y	Ν	Y	Ν	Ν	Ν
13.	Y	Ν	Ν	Y	Y	Y
14.	Y	Ν	Ν	Y	Ν	Ν
15.	Y	Ν	Ν	Ν	Y	Y
	•••					
32.	Ν	Ν	Ν	Ν	N	Ν

Table for abductive diagnosis



Using the table, we look for minimal sets of literals.

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For example, from line (5), we have that
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Ab(b1) \wedge Ab(b2) \wedge \neg Ab(a1) \wedge Ab(a2) \wedge Ab(o1)
```

entails the observations. However, lines (5), (7), (13) and (15) together lead us to a smaller set of literals (the first explanation below).

The explanations are

- 1. $Ab(b1) \land \neg Ab(a1) \land Ab(o1)$
- 2. $Ab(b1) \land \neg Ab(a1) \land \neg Ab(a2)$
- **3**. Ab(b2) $\land \neg Ab(a1) \land Ab(o1)$

Note: not all components are mentioned since for these settings, get the same observations whether or not they are working

but for this fault model only

Can narrow down diagnosis by looking at a number of different settings differential diagnosis

Diagnosis revisited

Abductive definition has limitations

- often only care about what is not working
- may not be able to characterize all possible failure modes
- want to prefer diagnoses that claim as few broken components as possible

Consistency-based diagnosis:

Assume KB uses the predicate ${\rm Ab}$ as before, but perhaps only characterizes the normal behaviour

e.g. And gate(x) $\land \neg Ab(x) \supset out(x) = and(in1(x), in2(x))$

Want a minimal set of components D, such that

 $\{Ab(c) \mid c \in D\} \cup \{\neg Ab(c) \mid c \notin D\}$

can use table as before with last column changed to "consistency"

is consistent with KB \cup Settings \cup Observations

In previous example, get 3 diagnoses: {b1}, {b2, a2} and {b2,o1}

Note: more complex to handle non-minimal diagnoses

Some complications

- 1. negative evidence
 - allow for missing observations

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e.g. ensure that \mathsf{KB} \cup \{\alpha\} \mid \neq \text{ fever}
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- 2. variables and quantification
 - same definition, modulo "simplicity", (but how to use Resolution?)
 - useful to handle open wffs also

KB $\cup \{x = 3\} \models P(x)$ handles WH-questions

- 3. probabilities
 - · not all simplest explanations are equally likely
 - also: replace (*Disease* \land ... \supset *Symptom*) by a probabilistic version
- 4. defaults
 - instead of requiring KB \cup { α } |= β , would prefer that given α , it is reasonable to believe β

e.g. being a bird explains being able to fly

Other applications

1. object recognition

what scene would account for image elements observed? what objects would account for collection of properties discovered?

2. plan recognition

what high-level goals of an agent would account for the actions observed?

3. hypothetical reasoning

instead of asking: what would I have to be told to believe β ? ask instead: what would I learn if I was told that α ?

Dual of explanation: want β such that

 $\begin{array}{l} \mathsf{KB} \cup \{\alpha\} \models \beta \\ \mathsf{KB} \not\models \beta \\ \mathsf{simplicity, parsimony} \\ \mathsf{using correct vocabulary} \end{array}$

can use the abduction procedure

• International Workshop on the Principles of Diagnosis

22nd International Workshop on Principles of Diagnosis



International Diagnostic Competition

Example Challenge Problem



Real World Diagnosis

- Key inferences
 - Fault detection
 - Fault isolation
 - Fault identification
 - Fault recovery
- Some challenges
 - Insufficient knowledge
 - Limited observability
 - Mixture of faults
 - Delayed effects
 - Fault Masking

- Prevalent techniques
 - Expert systems
 - Model based systems
 - Data driven systems
 - Stochastic systems
- Prevalent Algorithms
 - Mostly hybrid
 - Leverage model-based approaches
 - Consistency based inference

Suggested Reading

• Chapter 13 in Brachman & Levesque